

Approximate BDD Optimization

Rolf Drechsler^{1,2}

drechsler@uni-bremen.de

¹ Deutsches Forschungszentrum
für Künstliche Intelligenz (DFKI)
FB Cyber-Physical Systems

² University of Bremen
Group of Computer Architecture



Joint Work with



- Arun Chandrasekharan
- Saman Fröhlich
- Daniel Große
- Oliver Keszöcze
- Saeideh Shirinzadeh
- Mathias Soeken

Deutsche
Forschungsgemeinschaft

DFG

Introduction



- Approximate Computing (AC): consider applications like



robotics



search engines



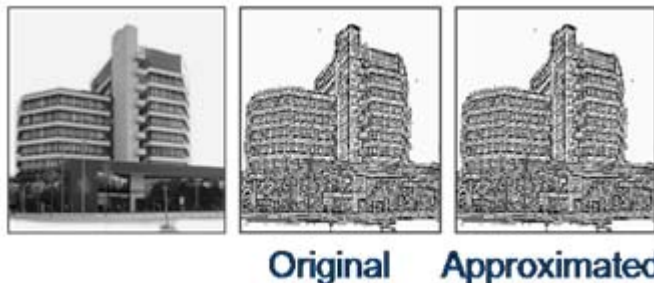
digital image processing

- What means “correct”?

Introduction



- Many real world applications tolerate inaccuracies
 - Media processing, recognition, data mining, etc.



- Usually due to *timing induced errors* (e.g. voltage-scaling) or *functional approximation*

- **AC trades off** inaccuracies for performance

Why BDDs?

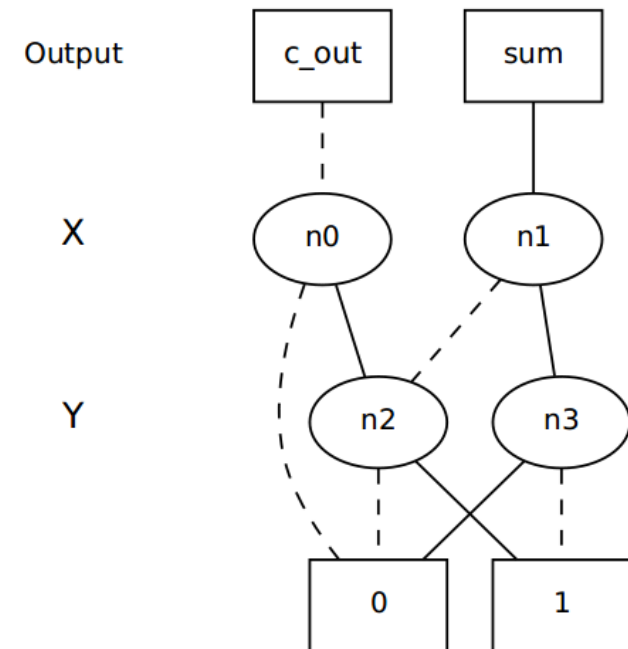


- Successfully applied in many fields
 - Test, verification, logic synthesis
- BDDs are well understood
 - Efficient algorithms
 - Lower/upper bounds on size of BDDs
- Allows deeper understanding of AC

Is it a new problem?



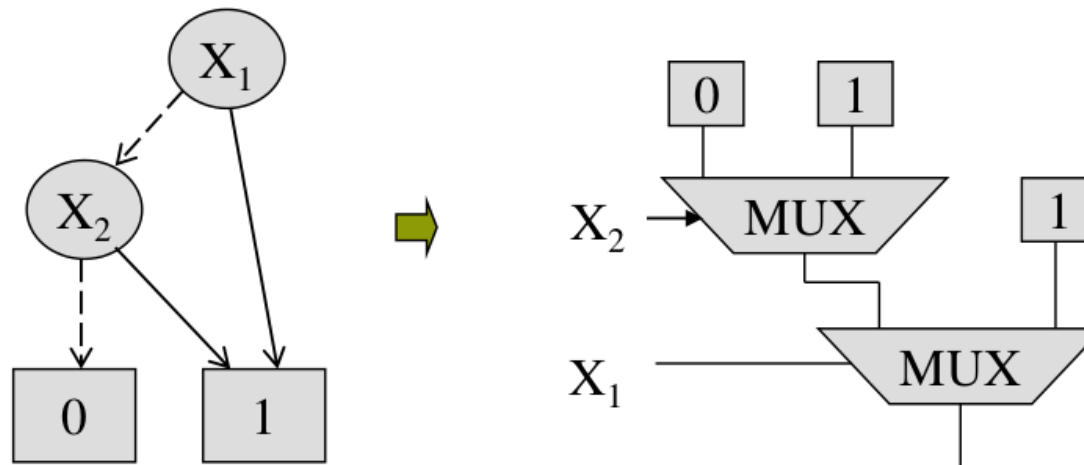
- Optimization of BDDs
 - Variable ordering
 - Assignment of **don't cares** (DCs)
 - AC: dynamically changing the DC set



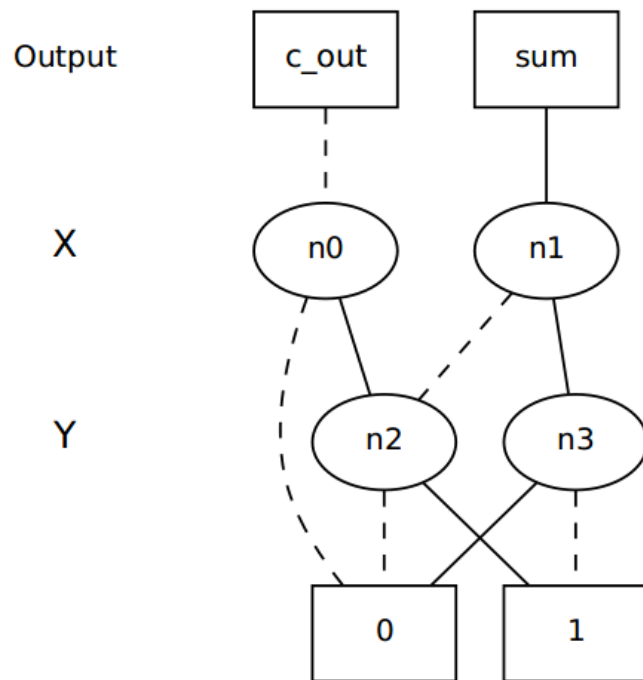
Reduction of Size



- Design styles based on direct mapping



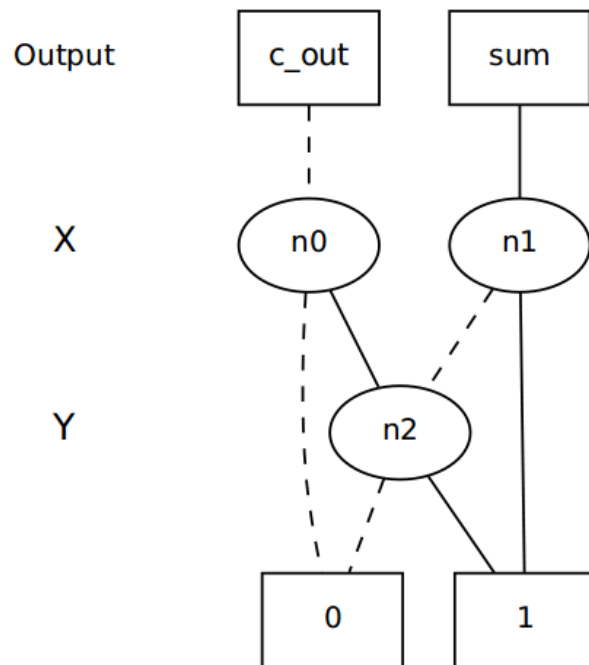
BDD for Half Adder



X	Y	Sum	C_out
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1

- Remove n3

Compacted BDD for Half Adder



X	Y	Sum	C_out
0	0	0	0
0	1	1	0
1	0	1	0
1	1	1	1

- Smaller hardware representation
- Error in the result needs to be tolerated

- Worst-Case:
 - $wc(f, \hat{f}) = \max_x \{ |\text{int}(f(x)) - \text{int}(\hat{f}(x))| \}$
- Average-Case:
 - $ac(f, \hat{f}) = \frac{\sum_x |\text{int}(f(x)) - \text{int}(\hat{f}(x))|}{2^n}$
- Error-Rate:
 - $er(f, \hat{f}) = \frac{\sum_x |f(x) \neq \hat{f}(x)|}{2^n}$
- ...

Evaluating Error Metrics

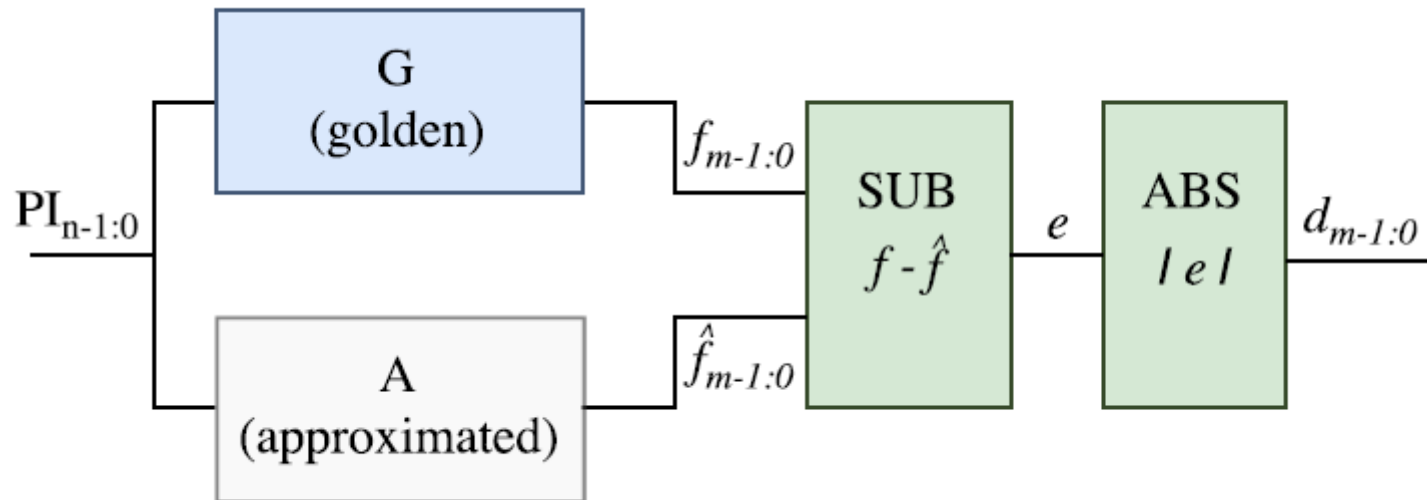


- $wc(f, \hat{f})$ is FP^{NP} – complete
 - Proof: $wc(f, \hat{f})$ is equivalent to LEXSAT
- $ac(f, \hat{f})$ is #P – complete
 - Proof: $ac(f, \hat{f})$ is equivalent to #SAT
- $er(f, \hat{f})$ is #P – complete
 - Proof: $ec(f, \hat{f})$ is equivalent to #SAT
- Similar complexities for other error-metrics

Evaluation using Mitters

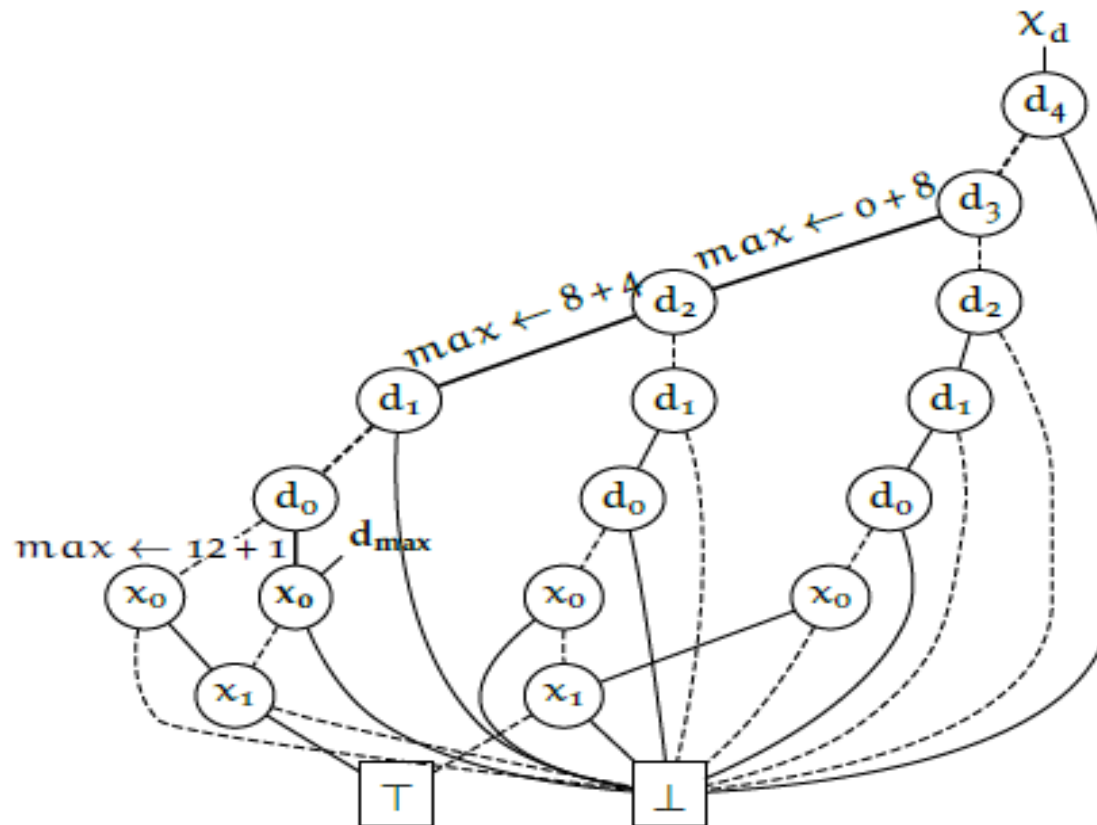


- Combinatorial circuits:



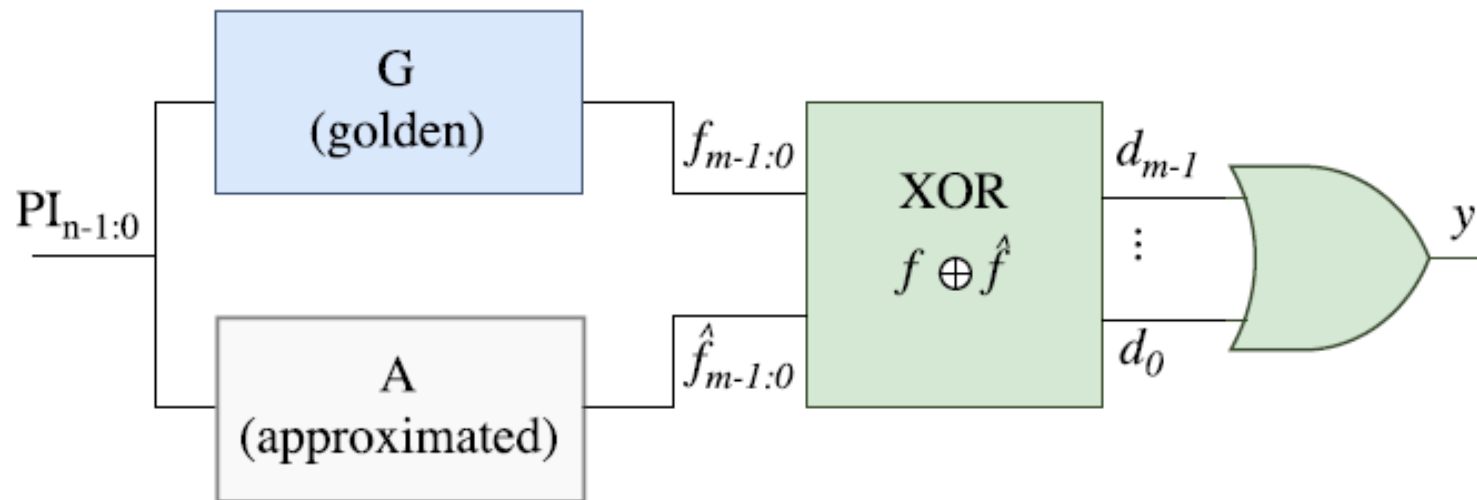
- Methodology:
 - Represent Miter as a BDD
 - Evaluate BDD in terms of the given error-metric

Worst-case Error



- similar algorithm exists for the average-case error

Error Rate

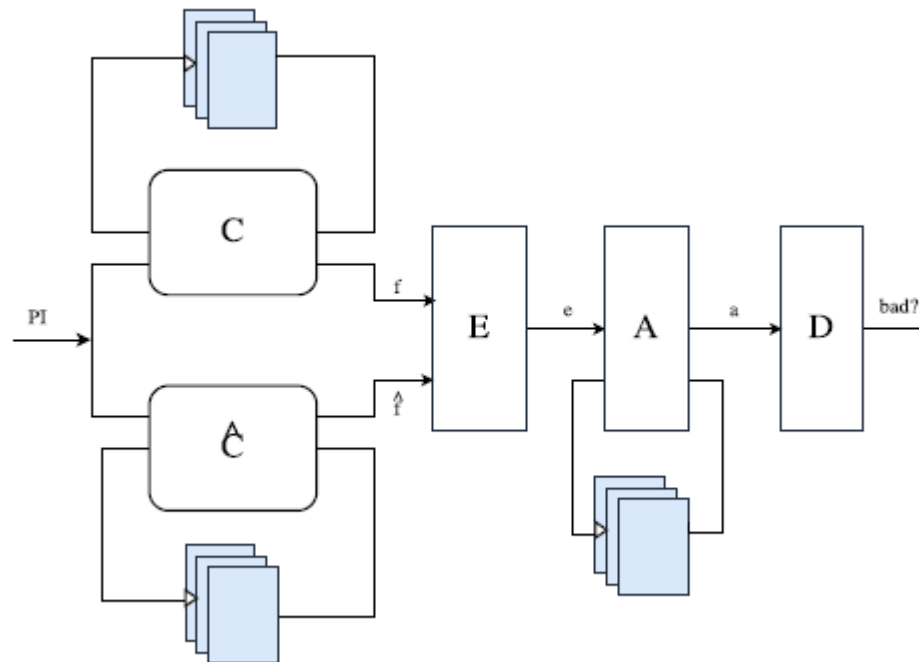


- Calculate ON-set of BDD-representation

Sequential Circuits



- E: error computation, A: accumulator, D: decision



- Use Bounded Model Checking (BMC) and Property Directed Reachability (PDR) techniques to determine error-metrics



- Directly on BDD:
 - Exact Minimization
 - Heuristic
- Evolutionary algorithms

Exact BDD Minimization



Error Bounded Exact BDD Minimization

Given:

$$f : \mathbb{B}^n \rightarrow \mathbb{B}$$
$$e \in \mathbb{N}$$

Task:

$$\min B(\hat{f}), \hat{f} : \mathbb{B}^n \rightarrow \mathbb{B}, \text{ s.t.}$$
$$\sum_{x \in \mathbb{B}^n} [\hat{f}(x) \neq f(x)] \leq e$$

Idea of Exact BDD Minimization



- Flip at most e bits in truth table of f
- Build a BDD representing all feasible functions
- Traverse BDD to find result

Experiment: Exact BDD Minimization



TABLE I
COMPUTATION TIME FOR BDD-SOLVER FOR FUNCTIONS WITH 3 VARIABLES

e	total BDD [s]	at BDD [s]	wt BDD [s]
1	0.018062	0.00007	0.00012
2	0.030604	0.00012	0.00018
3	0.021506	0.00008	0.00032

TABLE II
COMPUTATION TIME FOR BDD-SOLVER FOR FUNCTIONS WITH 4 VARIABLES

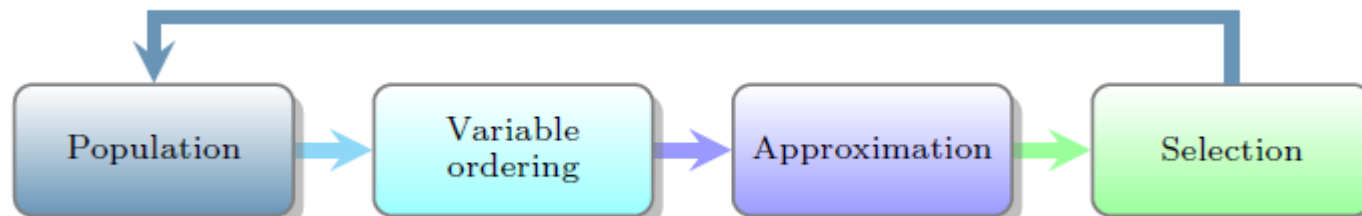
e	total BDD [s]	at BDD [s]	wt BDD [s]
1	21.28	0.00032	0.00102
2	100.07	0.00153	0.00260
3	444.73	0.00679	0.01300
4	1,289.70	0.01968	0.03000
5	2,798.10	0.04270	0.10200
6	3,948.30	0.06025	0.15000
7	2,392.00	0.03650	0.22000

- Does not scale well enough to be practical for large functions
 - Can be used to benchmark the quality of heuristics

Evolutionary Approach



- Based on evolutionary multi-objective optimization (MOO)
 - Variable reordering together with applying approximation operators



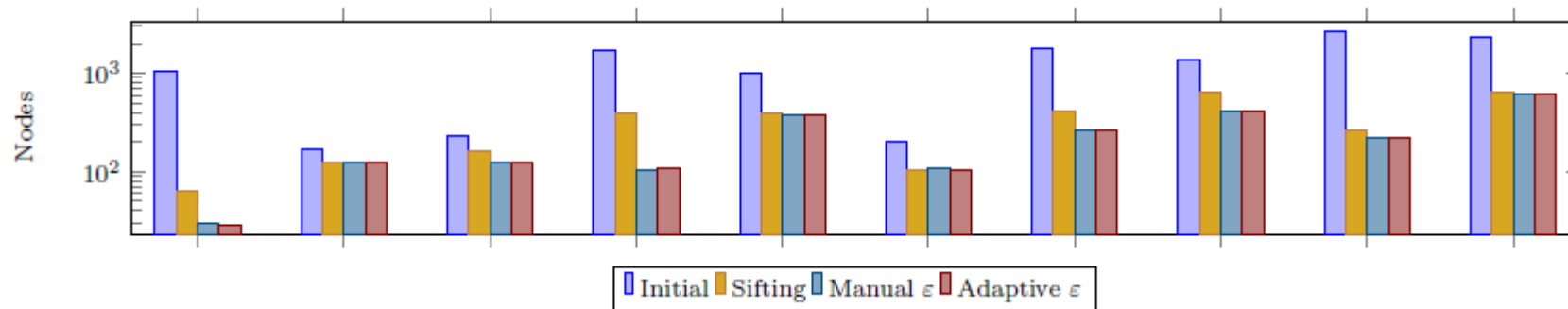
- Algorithm requires
 - Higher priority of BDD size
 - Error constraint satisfaction

- Evolutionary Multi-objective Optimization
 - $\min_x f(x) = (f_1(x), f_2(x), \dots, f_n(x))$
 - Pareto-front, Pareto-dominance
- Optimization goals are:
 - BDD size
 - Error

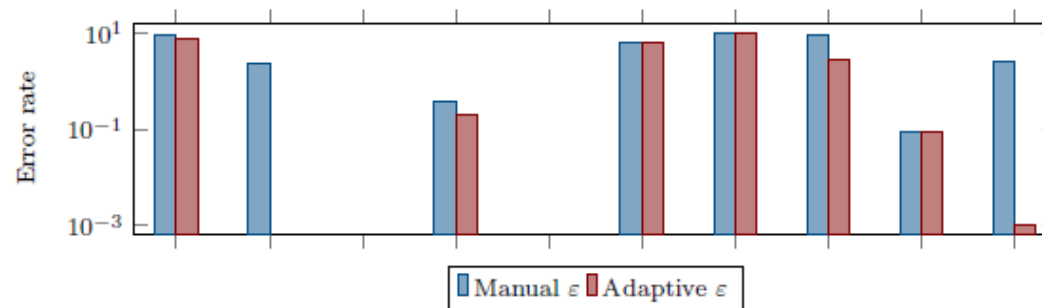
Evolutionary Approach: Experiments



- **68.02%** size reduction at a cost of 2.12% inaccuracy on ISCAS89
- **23.51%** size reduction in comparison with sifting



- **25.19%** error rate reduction by adaptive approach



References



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- BDDs for AC
- BDDs are intensively studied data structure
 - Theoretical analysis
 - Implementation and experiments
- Future work: more general structures

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